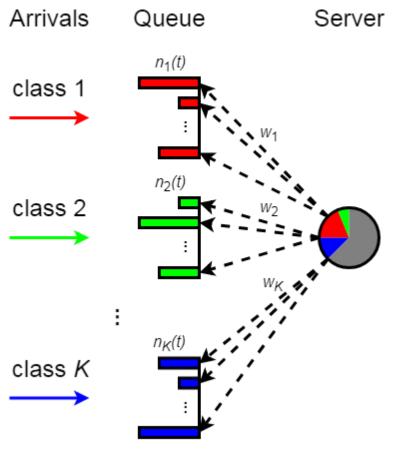
# Fluid Approximation of Closed Queueing Networks with Discriminatory Processor Sharing

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### **DPS** Discipline



Instantaneous fraction of the service capacity allotted to a class-k job:

$$\frac{w_k}{\sum_{l=1}^K w_l n_l(t)}$$

### **DPS** Queueing System

### **Related Work**

- Heavy-Traffic Regime:
  - A Closed Network with a Discriminatory Processor-Sharing Server [Mitra & Weiss, 1989]  $\frac{dx_k(t)}{dt} = \lambda_k (b_k - x_k(t)) - \frac{\lambda_k g_k w_k x_k(t)}{\sum_{l=1}^K w_l x_l(t)}$
- Moderately Heavy-Traffic Regime:
  - Asymptotic analysis of a large closed queueing network with discriminatory processor sharing [Morrison, 1991]

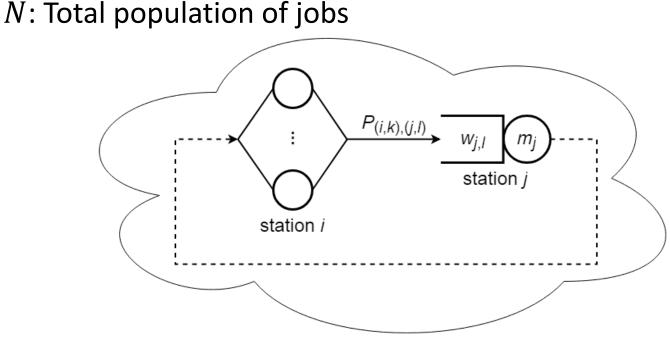
$$\sum_{l=1}^{K} \frac{b_k}{g_k} = 1 - \frac{a}{\sqrt{N}}$$

where a = O(1) as  $N \to \infty$ 

## **Model Definition**

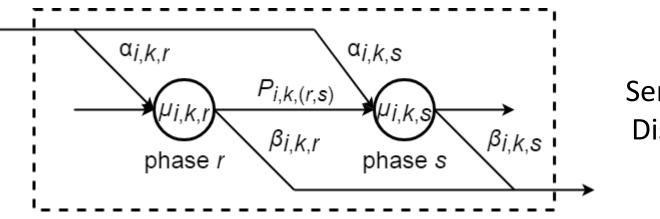
- *M*: Number of stations
- *K*: Number of classes

- $\mathcal{D}\colon \mathsf{Set} \ \mathsf{of} \ \mathsf{delay} \ \mathsf{stations}$
- $\mathcal{Q}$  : Set of queueing stations



### **Reference Model**

### **Model Definition**



Service Time Distribution

 Probability that a class-k job in phase r in station i proceeds to station j as a class-l job in phase s:

$$P_{(i,k,r),(j,l,s)} = \begin{cases} P_{i,k,(r,s)} + \beta_{i,k,r} P_{(i,k),(i,k)} \alpha_{i,k,s} & \text{if } i = j, k = l \\ \beta_{i,k,r} P_{(i,k),(j,l)} \alpha_{j,l,r} & \text{otherwise} \end{cases}$$

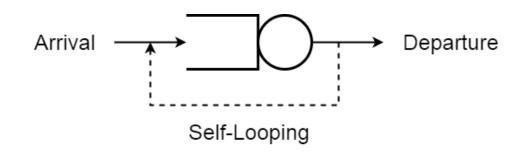
where

 $\beta_{i,k,r} = 1 - \sum_{s=1}^{R_{i,k}} P_{i,k,(r,s)}$ 

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### Fluid Approximation

- Standard Markovian Technique: State space explosion
- Fluid Approximation: Discontinuity of the ODE system
  - Preload a self-looping job into every queueing station:



 Assign the self-looping job in each queueing station i a service weight:

$$w_{i,0} = \frac{\sum_{k=1}^{K} w_{i,k}}{K}$$

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### Fluid Approximation

• Constraints on the service weight of the self-looping job:

$$\begin{cases} w_{i,0} \ge 0 & \text{if } \forall k \in \mathbb{N}_{\le K}^*, w_{i,k} \ge 0 \\ w_{i,0} = a & \text{if } \forall k \in \mathbb{N}_{\le K}^*, w_{i,k} = a \end{cases}$$

where  $a \in \mathbb{R}_{\geq 0}$  is arbitrary

 Assume the model to be a multi-chain QN, where no transient classes would possibly exist

### Fluid Approximation

• State of the model:  $\boldsymbol{n}(t) = \left(n_{i,k,r}(t)\right) \in \left\{\boldsymbol{n} \in \mathbb{N}^{R} : \sum_{i=1}^{M} \sum_{k=1}^{K} \sum_{r=1}^{R_{i,k}} n_{i,k,r} = N\right\}$ where

$$R = \sum_{i=1}^{M} \sum_{k=1}^{K} R_{i,k}$$

• Set of possible state changes in the model: L  $= \left\{ -\boldsymbol{e}_{i,k,r} + \boldsymbol{e}_{j,l,s} : i, j \in \mathbb{N}^*_{\leq M}; k, l \in \mathbb{N}^*_{\leq K}; r \in \mathbb{N}^*_{\leq R_{i,k}}; s \in \mathbb{N}^*_{\leq R_{j,l}} \right\}$ 

where  $e_{i,k,r}$  is a vector of all zeros except for a one in the (i,k,r)-th entry and  $e_{j,l,s}$  is similarly defined

### **Fluid Approximation**

• Transition rate of the model:

$$q_{\boldsymbol{n},\boldsymbol{n}+\boldsymbol{l}} = f(\boldsymbol{n},\boldsymbol{l})$$

where 
$$f : \mathbb{R}^{R} \times \mathbb{Z}^{R} \to \mathbb{R}_{\geq 0}$$
 is a function defined by  

$$f(\mathbf{x}, \mathbf{l}) = \begin{cases} \mu_{i,k,r} P_{(i,k,r),(j,l,s)} x_{i,k,r} & \text{if } \mathbf{l} \in L, i \in \mathcal{D} \\ \mu_{i,k,r} P_{(i,k,r),(j,l,s)} \frac{m_{i} w_{i,k} x_{i,k,r}}{\sum_{k'=1}^{K} w_{i,k'} \left(\sum_{r'=1}^{R_{i,k'}} x_{i,k',r'} + \frac{1}{K}\right)} & \text{if } \mathbf{l} \in L, i \in \mathcal{D} \\ 0 & \text{otherwise} \end{cases}$$

The function f(x, l) is continuous, bounded and Lipschitz continuous on any non-empty compact and convex subset of its global domain

$$D = \left\{ \boldsymbol{x} \in \mathbb{R}^{R} : \forall i \in \mathcal{Q}, \sum_{k'=1}^{K} w_{i,k'} \left( \sum_{r'=1}^{R_{i,k'}} x_{i,k',r'} + 1/K \right) \neq 0 \right\}$$

### Fluid Approximation

**Lemma 1.** Let  $\{n^{[v]}(t) : v \in \mathbb{N}^*\}$  be the sample paths of an infinite sequence of QN models whose v-th element is the original model scaled with a total population of vN jobs as well as  $vm_i$  servers and v self-looping jobs in each queueing station i. The sequence  $\{n^{[v]}(t) : v \in \mathbb{N}^*\}$  consistutes a density-dependent family of CTMCs.

### Fluid Approximation

**Lemma 2.** Define a vector field  $F : \mathbb{R}^R \to \mathbb{R}^R$  by  $F(x) = \sum_l lf(x, l)$ 

Specifically,

$$F_{i,k,r}(\boldsymbol{x})$$
  
=  $\sum_{j=1}^{M} \sum_{l=1}^{K} \sum_{s=1}^{R_{j,l}} \left( f(\boldsymbol{x}, -\boldsymbol{e}_{j,l,s} + \boldsymbol{e}_{i,k,r}) - f(\boldsymbol{x}, -\boldsymbol{e}_{i,k,r} + \boldsymbol{e}_{j,l,s}) \right)$ 

The initial value problem (IVP)

$$\begin{cases} \frac{d\boldsymbol{x}(t)}{dt} = \boldsymbol{F}(\boldsymbol{x}(t)) \\ \boldsymbol{x}(0) = \boldsymbol{n}(0) \end{cases}$$

has a unique solution x(t) in any finite time horizon T.

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### Fluid Approximation

Theorem 1. Given the same initial state

$$\frac{\boldsymbol{n}^{[\nu]}(0)}{\boldsymbol{v}} = \boldsymbol{n}(0)$$

the normalized sequence  $\{n^{[v]}(t)/v : v \in \mathbb{N}^*\}$  converges to the solution x(t) to the IVP

$$\begin{cases} \frac{d\boldsymbol{x}(t)}{dt} = \boldsymbol{F}(\boldsymbol{x}(t)) \\ \boldsymbol{x}(0) = \boldsymbol{n}(0) \end{cases}$$

in the sense that

$$\lim_{v \to \infty} \mathbb{P}\left(\sup_{t \in [0,T]} \left\| \frac{\boldsymbol{n}^{[\boldsymbol{v}]}(t)}{\boldsymbol{v}} - \boldsymbol{x}(t) \right\|_{1} > \delta\right) = 0 \quad \text{for } \delta \in \mathbb{R}_{>0}$$

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### Fluid Approximation

**Corollary 1.** Suppose that the condition of Theorem 1 is satisfied. The expectation of the normalized sequence  $\{n^{[v]}(t)/v : v \in \mathbb{N}^*\}$  converges to the solution x(t) to the IVP

$$\begin{cases} \frac{d\boldsymbol{x}(t)}{dt} = \boldsymbol{F}(\boldsymbol{x}(t)) \\ \boldsymbol{x}(0) = \boldsymbol{n}(0) \end{cases}$$

in any finite time horizon T:

$$\lim_{v \to \infty} \mathbb{E}\left(\frac{\boldsymbol{n}^{[\boldsymbol{v}]}(t)}{v}\right) = \boldsymbol{x}(t) \quad \text{for } t \in [0, T]$$

### **Steady-State Validation**

- A large set of 4914 model instances was obtained by setting
  - Topology: serial, parallel
  - Service time distribution: exponential (insensitive to phase-type distributions)
  - Number of stations: M = 3, 5, 9 (including a delay station)
  - Number of classes: K = 2, 4, 8
  - Total population of jobs: N = 128, 256, 512 ( $N_k = N/K$  without class switching)
  - Service rates: balanced, unbalanced (randomly generated with 30 seeds)
  - Service weights: balanced, unbalanced (randomly generated with 30 seeds)

### **Steady-State Validation**

• The proposed approach was compared with an AMVA algorithm for DPS presented in *Quantitative System Performance - Computer System Analysis Using Queueing Network Models* [Lazowska et al., 1984]:

$$R_{i,k}(\mathbf{N}) \approx \begin{cases} D_{i,k} & \text{if } i \in \mathcal{Q} \\ D_{i,k} \left( 1 + \frac{\sum_{l=1}^{K} w_{i,l} Q_{i,l}(\mathbf{N} - \mathbf{e}_k)}{w_{i,k}} \right) & \text{otherwise} \end{cases}$$

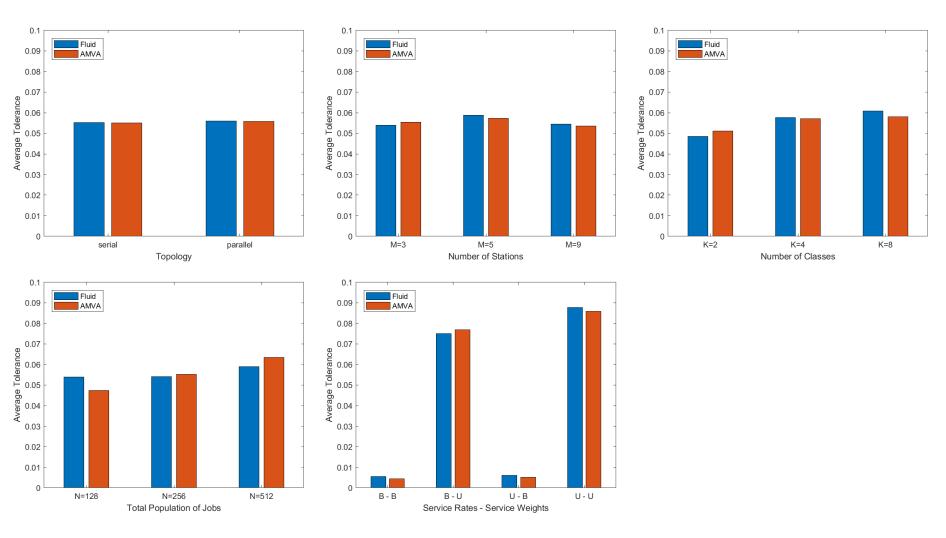
 The results were evaluated against discrete-event simulation using a tolerance metric found in *Linearizer - A Heuristic Algorithm for Queueing Network Models of Computing Systems* [Chandy & Neuse, 1982]:

$$\Gamma olerance = \max_{\substack{i \in \mathbb{N}^* \leq M \in \mathbb{N}^* \\ \text{Discriminatory Processor Sharing}}} \frac{\left|Q_{i,k}^{C} - Q_{i,k}^{R}\right|}{\mathbb{Q}_{i,k}^{C} - Q_{i,k}^{R}\right|}$$

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### **Steady-State Validation**

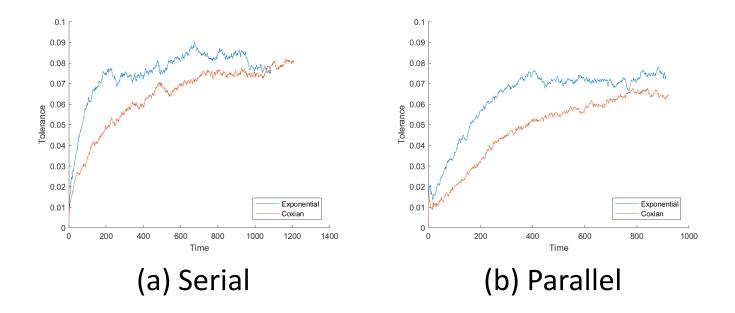


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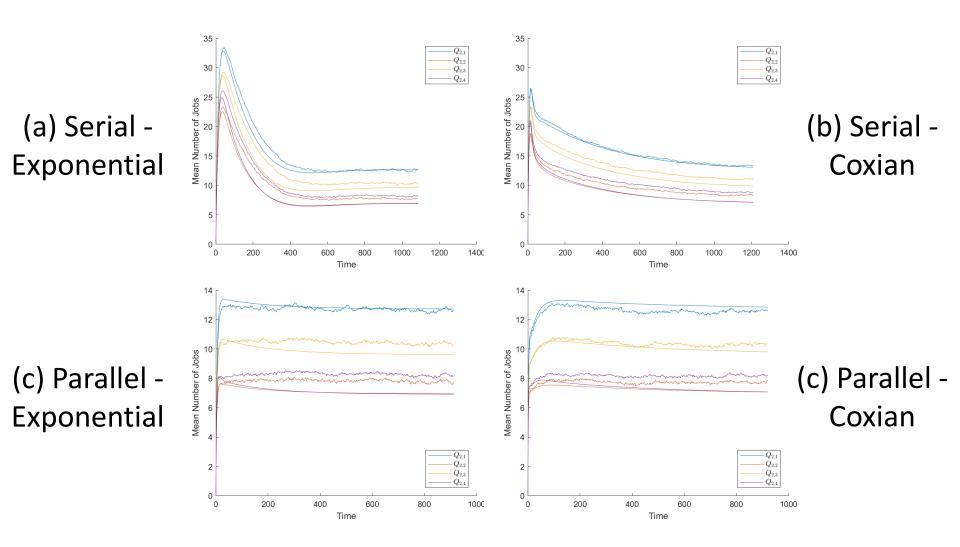
### **Transient Validation**

- 4 typical model instances were obtained by setting
  - Topology: serial, parallel
  - Service time distribution: exponential, 2-phase coxian with a CV of 2 (sensitive to phase-type distributions)
  - Number of stations: M = 5 (including a delay station)
  - Number of classes: K = 4
  - Total population of jobs: N = 256 ( $N_k = N/K$  without class switching)
  - Service rates: unbalanced (randomly generated with 1 seed)
  - Service weights: unbalanced (randomly generated with 1 seed)

### **Transient Validation**



### **Transient Validation**



Fluid Approximation of Closed Queueing Networks with Discriminatory Processor Sharing

### Applications

- Dynamic Analysis of Random Environments:
  - Blending randomness in closed queueing network models [Casale, Tribastone & Harrison, 2014]
- Approximation of Response Time Distributions:
  - LINE: Evaluating Software Applications in Unreliable Environments [Casale, Tribastone & Harrison, 2014]
  - Not accurate for non-exponential service times
  - Not applicable at the system level

## **Concept of Chains**

- The routing probability matrix  $[P_{(i,k),(j,l)}] \in \mathbb{R}^{MK \times MK}$  can be considered as defining a DTMC whose state is a station-class pair (i, k)
- For a multi-chain QN, the routing DTMC is assumed to be decomposable into ergodic subchains
- If we represent the routing DTMC as a weighted direct graph, then each chain manifests itself as a strongly connected subgraph
- Because all the chains are completely isolated from one another, it is impossible for a job to switch between different chains

### Response Time CDF

- $\mathcal{V}_q$ : Set of station-class pairs in chain q
- $\mathcal{E}_q$ : Set of routes connecting station-class pairs in chain q
- Population of jobs in chain q:

$$N_q^* = \sum_{(i,k)\in \mathcal{V}_q} \sum_{r=1}^{R_{i,k}} n_{i,k,r}$$

- Set of classes that may be visited a chain-q job at station i:  $C_{i,q} = \{k : (i,k) \in \mathcal{V}_q\}$
- Set of classes that are visited by chain-q jobs across all the stations:

$$\mathcal{C}_q = \bigcup_{i=1}^M \mathcal{C}_{i,q}$$

### **Response Time CDF**

- Evaluate the mean service time  $S_{i,k}$  of a class-k job at station i by applying the moment formula for phase-type distributions
- Modify the service time distribution of class k at station i so that any class-k arrival at the station starts in an additional phase 0 with probability  $\alpha_{i,k,0} = 1$  and shifts from that phase to phase r with probability  $P_{i,k,(0,r)} = \alpha_{i,k,r}$ 
  - Phase (i, k, 0) therefore gathers up all class-k jobs having just arrived at station i
  - To limit the side effects of the modification, the service rate of phase (i, k, 0) to  $\mu_{i,k,r} = 1/\eta S_{i,k}$ , where  $\eta$  is a small positive constant

### Response Time CDF

- Build a fluid ODE system that incorporates the modified service time distribution, and compute its steady-state solution  $\widetilde{x}' \in \mathbb{R}^{R'}$  under the initial condition  $x'(0) \in \mathbb{R}^{R'}$ , where R' = R + 1
  - The initial condition x'(0) derives from the given initial condition x(0) by inserting zero entries
- Create an auxiliary class k' for tagging class-k jobs
  - Class k' has the same service time distribution and the same service weight as class k at station i
  - The probability of a class-k' job at station i being routed to station j as a class-l job is set as  $P_{(i,k'),(j,l)} = P_{(i,k),(j,l)}$

• A tagged class-k job is thereby restored immediately 04/12/2018 after depart ure from station version Networks with Discriminatory Processor Sharing 24

### **Response Time CDF**

- Build another fluid ODE system that incorporates the auxiliary class, and compute its transient solution  $x''(t) \in \mathbb{R}^{R''}$  under the initial condition  $x''(0) \in \mathbb{R}^{R''}$ , where  $R'' = R + R_{i,k} + 2$ 
  - The initial condition x''(0) arises from the resultant steady-state solution  $\tilde{x}'$  by inserting zero entries except that  $x''_{i,k,0}(0) = 0$  and  $x''_{i,k',0}(0) = \tilde{x}'_{i,k,0}$
- Response time CDF of class k at station i:

$$\Phi_{i,k}(t) \approx 1 - \frac{\sum_{r=0}^{R_{i,k'}} x_{i,k',0}^{\prime\prime}(t)}{x_{i,k',0}^{\prime\prime}(0)}$$

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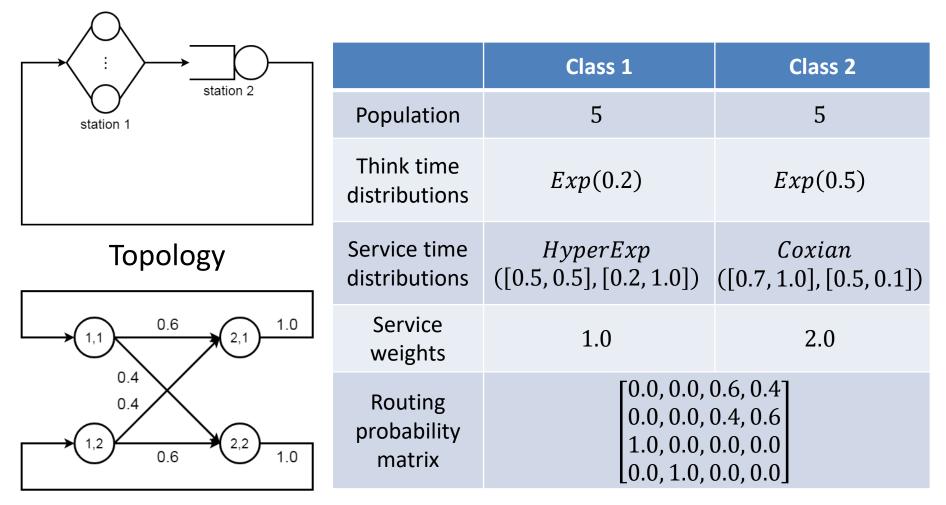
### System Response Time CDF

- Add an end phase in the service time distribution of each class k in the set  $\mathcal{C}_{D,q}$  at station D to tag chain-q jobs that are about to depart from the reference station and restore them upon next arrival.
- System response time CDF of chain q with respect to station
   D:

$$\Phi_{i,k}(t)$$

$$\approx 1 - \frac{\sum_{k \in \mathcal{C}_{D,q}} x_{D,k',R_{D,k'}+1}^{\prime\prime}(t) + \sum_{(i,k) \in \mathcal{V}_{q,i\neq D}} \sum_{r=1}^{R_{i,k'}} x_{i,k',r}^{\prime\prime}(t)}{\sum_{k \in \mathcal{C}_{D,q}} x_{D,k',R_{D,k'}+1}^{\prime\prime}(0)}$$

### A Simple Example



### Parameters

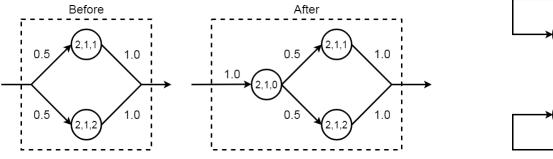
**Routing CTMC** 

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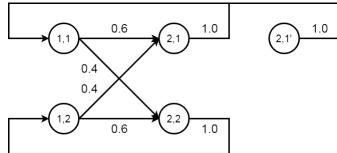
**Discriminatory Processor Sharing** 

### A Simple Example

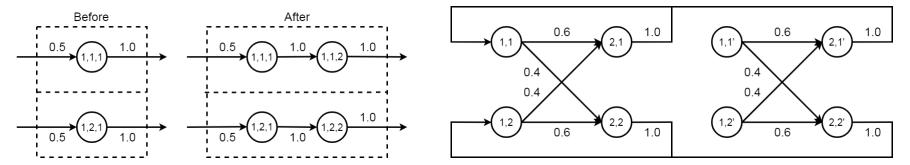
For Response Time CDF of Class 1 at Queueing Station



Service Time Distribution



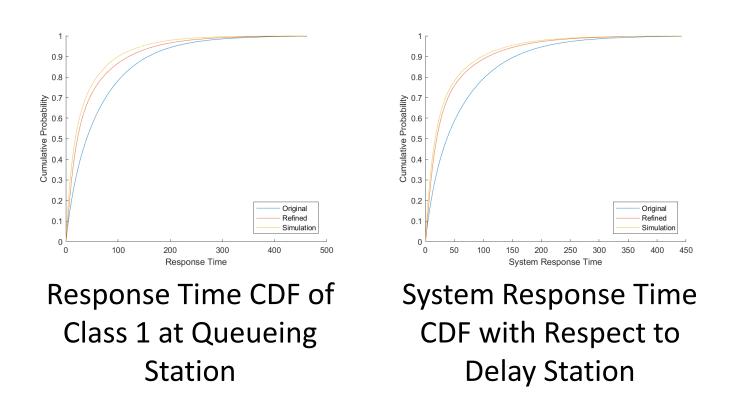
**Routing CTMC** 



For System Response Time CDF with Respect to Delay Station

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### A Simple Example



### Conclusion

- A fluid approach is proposed for approximating the transient and steady-state behavior of closed QNs with DPS
- Our reference model is defined generically, featuring
  - Arbitrary topology
  - Phase-type service times
  - Class switching
- The proposed approach has been verified against discreteevent simulation for both transient and steady-state analysis
- A refined method and its extension are introduced for approximating response time distributions at the station and system levels

Thank you!